

Almost Bayesian Assignments and Conditional Independence

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Outline of the Lecture

- 1 History and motivation
 - History
 - Motivation
 - Generating sequences
- 2 Belief function models
 - Set notation
 - Compositional models
 - Almost Bayesian basic assignments
- 3 Conditional independence
 - Conditional noninteractivity
 - Conditional independence

History

- 1986 Jiroušek Radim, Perez Albert: Graph-aided Knowledge Integration in Expert System INES. Proceedings IPMU'86.
- 1997 Jiroušek Radim: Composition of probability measures on finite spaces. Proceedings UAI'97.
- 1998 Vejnarová Jirina: Composition of possibility measures on finite spaces: Preliminary results. Proceedings IPMU'98.
- 2007 Jiroušek Radim, Vejnarová Jirina, Daniel Milan: Compositional models for belief functions. Proceedings ISIPTA'07.

Probabilistic operator of composition

For $\kappa_1(x_K)$ and $\kappa_2(x_L)$ defined on \mathbf{X}_K and \mathbf{X}_L , respectively, such that $\kappa_1^{\downarrow K \cap L} \ll \kappa_2^{\downarrow K \cap L}$, which means that

$$\forall x \in \mathbf{X}_{K \cap L} \quad (\kappa_2^{\downarrow K \cap L}(x) = 0 \implies \kappa_1^{\downarrow K \cap L}(x) = 0);$$

their composition is defined for all $x \in \mathbf{X}_{K \cup L}$

$$(\kappa_1 \triangleright \kappa_2)(x) = \frac{\kappa_1(x^{\downarrow K})\kappa_2(x^{\downarrow L})}{\kappa_2^{\downarrow K \cap L}(x^{\downarrow K \cap L})} = \kappa_1(x^{\downarrow K})\kappa_2(x^{\downarrow L \setminus K} | x^{\downarrow L \cap K}).$$

Basic properties of the operator of composition

- $\kappa_1(x_K) \triangleright \kappa_2(x_L) = (\kappa_1 \triangleright \kappa_2)(x_{K \cup L})$;
- $(\kappa_1(x_K) \triangleright \kappa_2(x_L)) \downarrow^K = \kappa_1(x_K)$;
- operator is neither commutative nor associative;
- $\kappa_1 \triangleright \kappa_2 = \kappa_2 \triangleright \kappa_1 \iff \kappa_1 \downarrow^{K \cap L} = \kappa_2 \downarrow^{K \cap L}$;
- $X_{K \setminus L} \perp\!\!\!\perp X_{L \setminus K} | X_{K \cap L} [\kappa_1 \triangleright \kappa_2]$;
- ...;
- $X_I \perp\!\!\!\perp X_J | X_K [\kappa] \iff \kappa(x \downarrow^{I \cup J \cup K}) = \kappa(x \downarrow^{I \cup K}) \triangleright \kappa(x \downarrow^{J \cup K})$.

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Generating sequences

$$\kappa_1(x_{K_1}), \kappa_2(x_{K_2}), \dots, \kappa_n(x_{K_n})$$

Definition

Compositional Model:

$$\begin{aligned} \kappa_1(x_{K_1}) \triangleright \kappa_2(x_{K_2}) \triangleright \dots \triangleright \kappa_n(x_{K_n}) \\ = \left(\dots \left(\kappa_1(x_{K_1}) \triangleright \kappa_2(x_{K_2}) \right) \triangleright \dots \triangleright \kappa_n(x_{K_n}) \right) \end{aligned}$$

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Perfect sequences

Definition

Generating sequence $\kappa_1(x_{K_1}), \kappa_2(x_{K_2}), \dots, \kappa_n(x_{K_n})$ is **perfect** if

$$\kappa_1 \triangleright \kappa_2 = \kappa_2 \triangleright \kappa_1,$$

$$(\kappa_1 \triangleright \kappa_2) \triangleright \kappa_3 = \kappa_3 \triangleright (\kappa_1 \triangleright \kappa_2),$$

$$(\kappa_1 \triangleright \kappa_2 \triangleright \kappa_3) \triangleright \kappa_4 = \kappa_4 \triangleright (\kappa_1 \triangleright \kappa_2 \triangleright \kappa_3),$$

...

$$(\kappa_1 \triangleright \dots \triangleright \kappa_{n-1}) \triangleright \kappa_n = \kappa_n \triangleright (\kappa_1 \triangleright \dots \triangleright \kappa_{n-1}).$$

Perfect sequences

Characterization Theorem:

Generating sequence $\kappa_1, \kappa_2, \dots, \kappa_n$ is perfect iff all κ_j are marginal distributions of $\kappa_1 \triangleright \kappa_2 \triangleright \dots \triangleright \kappa_n$.

Comparison with Bayesian networks

Both compositional models and Bayesian networks represent the same class of distributions

Pros

- It does not need help of Graph Theory
- Perfect sequence models are computationally more efficient

Cons

- Compositional models are not always defined
- It is more difficult to construct perfect sequence models

Operator of composition was defined also in Possibility theory (1998) and D-S theory of evidence (2007).

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Set notation

Let $K \subset L \subseteq N$ and $x \in \mathbf{X}_L$. $x^{\downarrow K}$ denotes a *projection* of x into \mathbf{X}_K .

Analogously, for and $A \subset \mathbf{X}_L$, $A^{\downarrow K}$ denotes a *projection* of A into \mathbf{X}_K :

$$A^{\downarrow K} = \{y \in \mathbf{X}_K \mid \exists x \in A : y = x^{\downarrow K}\}.$$

Important: we do not exclude $K = \emptyset$. In this case $A^{\downarrow \emptyset} = \emptyset$.

A *join* of two sets $A \subseteq \mathbf{X}_K$ and $B \subseteq \mathbf{X}_M$ is the set

$$A \otimes B = \{x \in \mathbf{X}_{K \cup M} : x^{\downarrow K} \in A \ \& \ x^{\downarrow M} \in B\}.$$

Operator of composition

For basic assignments m_1 on \mathbf{X}_K and m_2 on \mathbf{X}_L a *composition* $m_1 \triangleright m_2$ is defined for all $C \subseteq \mathbf{X}_{K \cup L}$ by one of the following expressions:

[a] if $m_2^{\downarrow K \cap L}(C^{\downarrow K \cap L}) > 0$ and $C = C^{\downarrow K} \otimes C^{\downarrow L}$ then

$$(m_1 \triangleright m_2)(C) = \frac{m_1(C^{\downarrow K}) \cdot m_2(C^{\downarrow L})}{m_2^{\downarrow K \cap L}(C^{\downarrow K \cap L})};$$

[b] if $m_2^{\downarrow K \cap L}(C^{\downarrow K \cap L}) = 0$ and $C = C^{\downarrow K} \times \mathbf{X}_{L \setminus K}$ then

$$(m_1 \triangleright m_2)(C) = m_1(C^{\downarrow K});$$

[c] in all other cases

$$(m_1 \triangleright m_2)(C) = 0.$$

Basic properties

For ba's m_1 on \mathbf{X}_K and m_2 on \mathbf{X}_L :

- $m_1 \triangleright m_2$ is ba on $\mathbf{X}_{K \cup L}$;
- m_1 is marginal of $m_1 \triangleright m_2$;
- operator is neither commutative nor associative;
-;
- for all $A \subseteq \mathbf{X}^{\downarrow K \cup L}$, $A \neq A^{\downarrow K} \otimes A^{\downarrow L} \implies (m_1 \triangleright m_2)(A) = 0$;

E.g. for binary case $|\mathbf{X}_{\{1,2,3\}}| = 2^8 - 1 = 255$

$$|\{A : A = A^{\downarrow \{1,2\}} \otimes A^{\downarrow \{2,3\}}\}| = 99.$$

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E.g. for binary case $|\mathbf{X}_{\{1,2,3\}}| = 2^3 - 1 = 7$

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Almost Bayesian basic assignments

Definition

Basic assignment is said to be **almost Bayesian** if it is

- *cylindrical* - all its focal elements C are point-cylinders ($C = C^{\downarrow L} \times \mathbf{X}_{K \setminus L}$ for $|C^{\downarrow L}| \leq 1$); and
- *sparse (quasi-Bayesian)*- all its focal elements are pairwise disjoint.

Assertion

Any compositional model assembled from Bayesian basic assignments is almost Bayesian.

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Conditional noninteractivity for belief functions

There are several definitions of this notion:

The most frequent is that used by M. Studený, P. Shenoy, Ben Yaghlane et al.:

Definition

Variables X_I and X_J are *conditionally non-interactive* given variables X_K ($X_I \perp\!\!\!\perp_{[m]} X_J | X_K$) if for all $A \subseteq \mathbf{X}_N$

$$\begin{aligned} Com_{m \downarrow I \cup J \cup K}(A \downarrow I \cup J \cup K) &\cdot Com_{m \downarrow K}(A \downarrow K) \\ &= Com_{m \downarrow I \cup K}(A \downarrow I \cup K) \cdot Com_{m \downarrow J \cup K}(A \downarrow J \cup K). \end{aligned}$$

$$Com_m(A) = \sum_{B \supseteq A} m(B).$$

Conditional independence for belief functions

In this case, however:

$$X_I \perp\!\!\!\perp_{[m]} X_J | X_K \not\iff m^{\downarrow I \cup J | K} = m^{\downarrow I | K} \triangleright m^{\downarrow J | K}$$

If we substitute conditional irrelevance by factorization we get conditional independence relation with the following properties:

- its restriction to Bayesian basic assignment corresponds to probabilistic conditional independence relation;
- it meets all the semigraphoid axioms (symmetry and Block independence lemma);
- it does not suffer from the conditional irrelevance imperfectness: it is consistent with marginalization.

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Consistency with marginalization

Having two projective ba's m_1 on \mathbf{X}_K and m_2 on \mathbf{X}_M ,
 (i.e. $m_1 \downarrow^{K \cap M} = m_2 \downarrow^{K \cap M}$) does there exist a ba m such that:

- m_1 and m_2 are marginal assignments of m ;
- $X_{K \setminus M} \perp\!\!\!\perp_{[m]} X_{M \setminus K} | X_{K \cap M}$?

The solution is simple:

$$m = m_1 \triangleright m_2.$$

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THANK YOU FOR YOUR ATTENTION